



On the Overhead Factors of Optimized LDGM– Staircase and LDGM–Triangle FEC Codes

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Overview

- Introduction
- DF vs. LDPC FEC Schemes
- Overhead Factor
- LDGM-type FEC Codes
- Simulation Results
- Conclusions



Introduction

- Large Scale Content Distribution systems:
 - Increasingly popular Internet application
 - To distribute a large file from one source to many nodes in the network
- Possible Scenarios:
 - Uncoded block transmission over Overlay (P2P) Networks (e.g. Bittorrent)
 - Coding allowed only at the source node
 - Digital Fountain (DF) approach
 - LDPC-based FEC approach
 - (Network) Coding allowed at intermediate nodes (e.g. Microsoft Avalanche)



DF vs. LDPC FEC approach

- DF approach:
 - Chop original file into K equal-length packets
 - Produce infinite amount of encoded packets by randomly XOR-ing subsets of original packets
 - Critical Part: Design of Degree Distributions
 - Any slightly more than K encoded packets enough for successful file recovery at the destinations
 - Universally nearly optimal



DF vs. LDPC FEC approach

- LDPC FEC approach
 - Finite approximation of DF
 - Select appropriate rate $R = K/N$ LDPC code to handle worst channel
 - Encode original K packets into length N codeword
 - Permute and send encoded packets
 - “Greedy” BP decoder at receiver
 - Hope for enough encoded packets to arrive at the destinations (Minimize Overhead Factor)



LDPC Code Design

- Design of LDPC Codes for BEC Channel
 - Theoretically mature
 - Asymptotically solved
 - Use capacity–approaching optimized (λ, ρ) –irregular LDPC ensemble
 - Finite–length design
 - From a given ensemble of LDPC codes select one with the size of the Minimal Stopping Set maximized
 - Unsolved.



Overhead Factor

- Overhead Factor F
 - Performance metric of interest in Content Distribution applications
 - Quantifies average amount of excessive encoded packets needed for the file reconstruction
 - $K' = (1 + F) \cdot K$
 - $K' > K$ is average number of encoded packets needed for file recovery
 - We got used to BER Curves, Waterfall and Error Floor Regions, DE Thresholds... Any Connections?



Asymptotic Overhead Factor

- Asymptotic Overhead Factor F_∞
 - $\epsilon^{Th}(\lambda, \rho)$: DE Threshold for (λ, ρ) LDPC ensemble
 - Ensemble rate $R = 1 - \epsilon^{Th}(\lambda, \rho)$ over $\text{BEC}(\epsilon^{Th}(\lambda, \rho))$
 - As $N \rightarrow \infty$ number of erased packets sharply concentrate around $N \cdot \epsilon^{Th}(\lambda, \rho)$ (Typicality)
 - Since we're on an ensemble capacity limit $N \cdot (1 - \epsilon^{Th}(\lambda, \rho))$ received packets is enough to drive error probability $\rightarrow 0$
 - Equalizing this with $(1 + F_\infty) \cdot K$
 - $F_\infty = \frac{1 - \epsilon^{Th}}{R} - 1$



Overhead Factor on Finite Lengths

- Finite-Length Overhead Factor
 - Should be related to Error-Floor performance?
 - Optimizing Error-Floor behavior:
Maximizing Minimal Stopping Sets
 - Applying State-of-the-art design algorithms that optimizes Error-Floor behavior and check what happens with an Overhead Factor at short-to-medium code lengths



Large Content Distribution Systems

- Standardization work
 - DF Raptor Codes as a core mechanism for
 - 3GPP MBMS (Multimedia Broadcast/Multicast System)
 - DVB-H IP Content Delivery
 - Possibilities for Patent-Free Codecs are left open
 - LDGM-type design codes are suggested
 - V. Roca and C. Neumann “Low Density Parity Check (LDPC) Forward Error Correction”, June 2005. IETF Draft Work in Progress: <draft-roca-rmt-ldpc-00.txt>

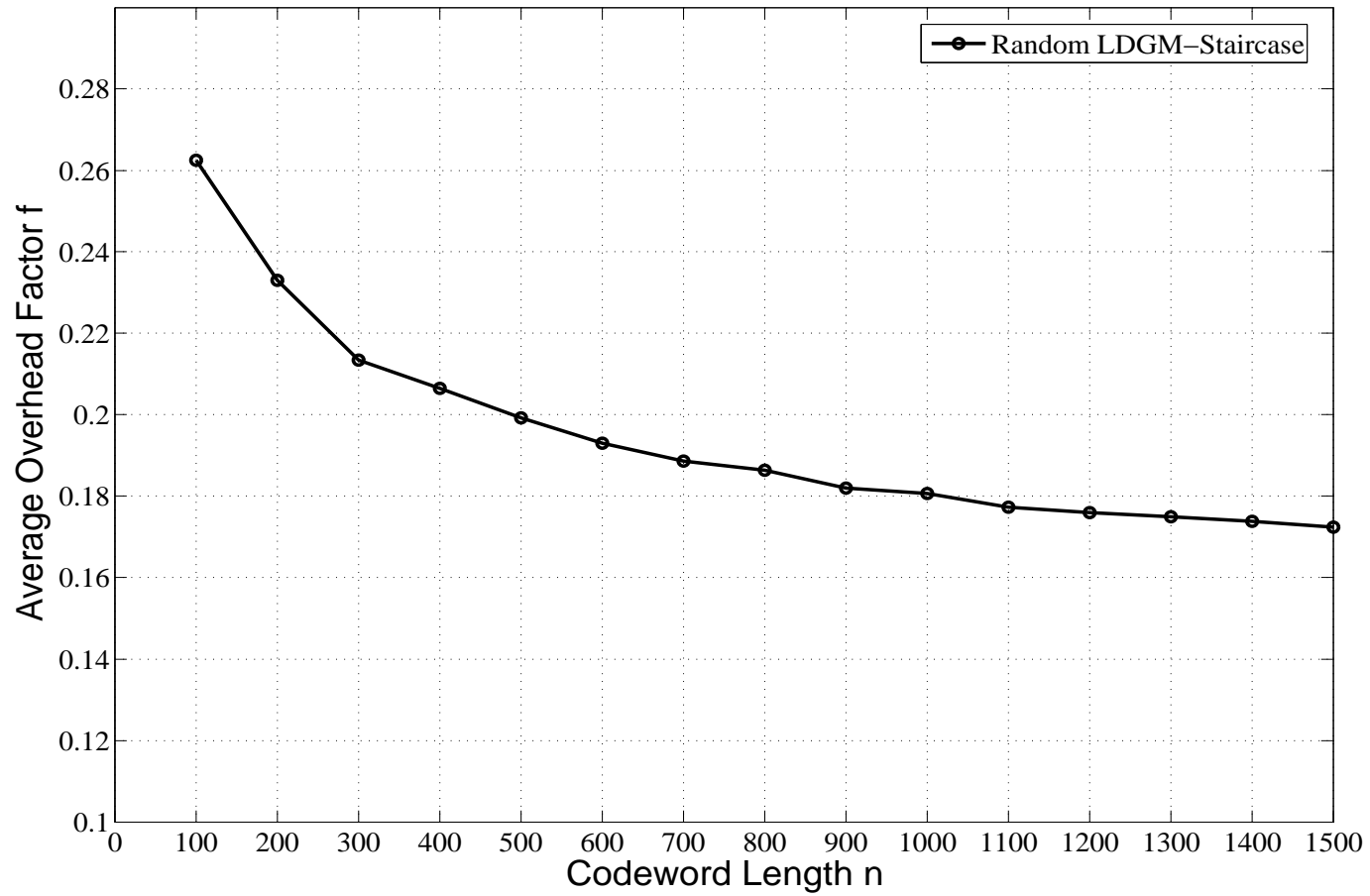


Simulation Results

- Simulation Setting
 - Overhead Factor F is examined by MC simulation for $R = \frac{1}{2}$ left-quasi regular LDGM–Staircase codes of increasing length from 100–1500 with step 100
 - For each length 1000 experiments with random selection of H and random permutation of codeword symbols before transmission
 - All-zero codeword assumption
 - Based on all generated H matrices, average degree distribution pair (λ, ρ) and corresponding $\epsilon^{Th}(\lambda, \rho)$ is calculated

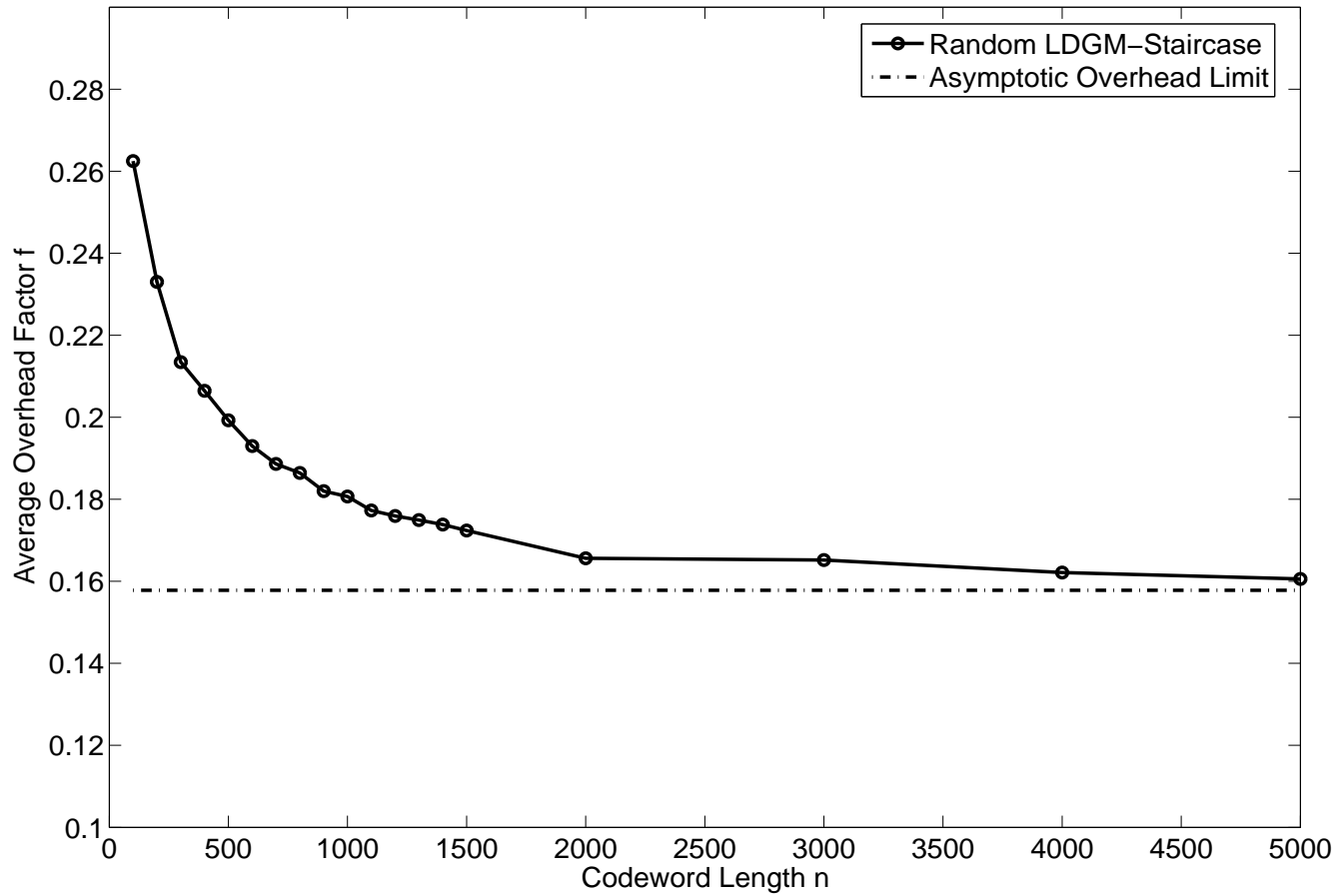


Random LDGM Staircase Design

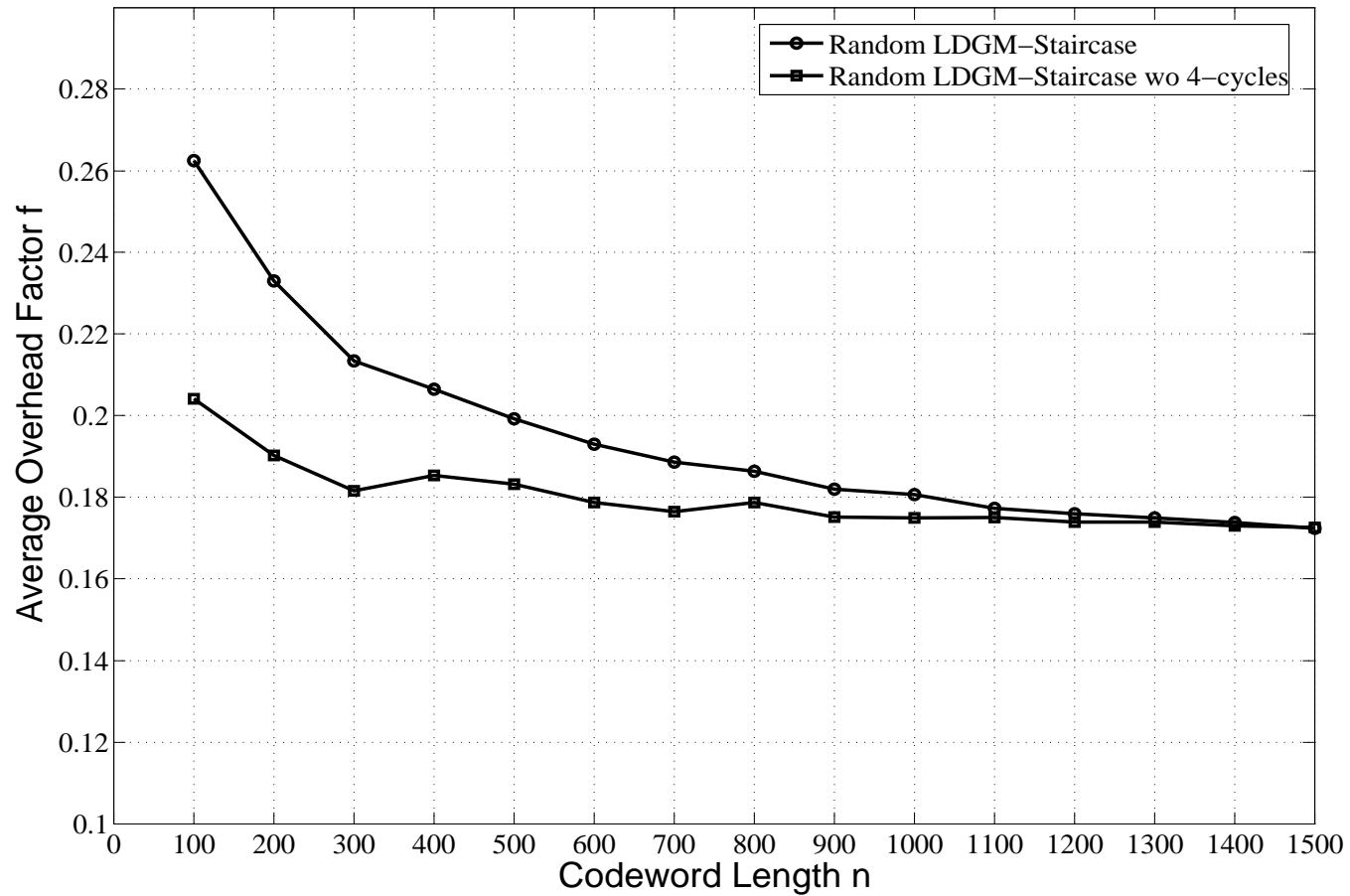




Asymptotic Overhead Factor



Random LDGM Staircase Design without 4-cycles



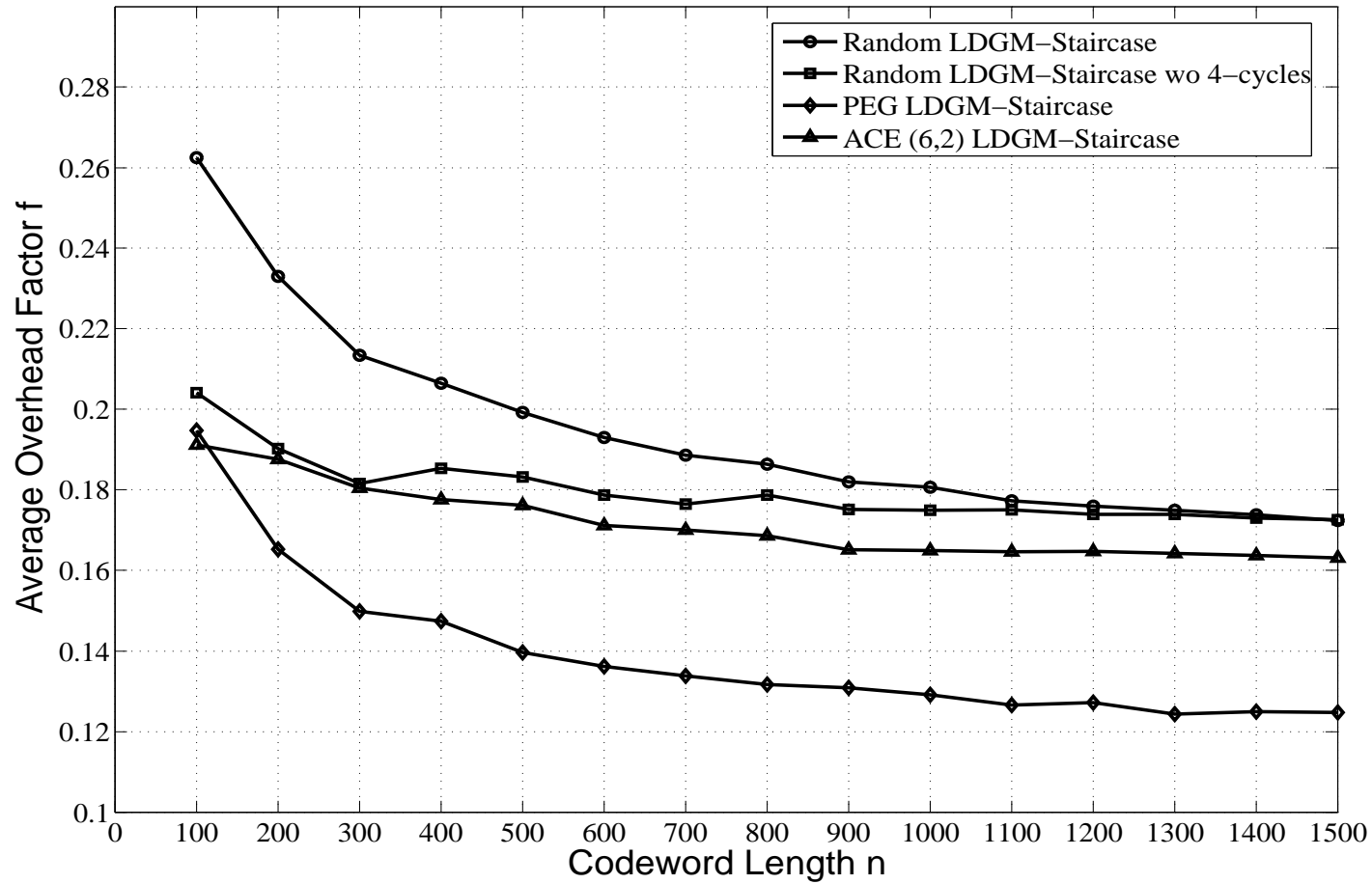


Improved LDGM–Staircase Design

- Two well-known algorithms
 - Progressive Edge Growth (PEG) Algorithm
 - [Hu, Eleftheriou, Arnold (2002)]
 - Heuristic algorithm for large girth graph design
 - ACE Constraining Algorithm
 - [Tian, Jones, Villasenor, Wesel (2004)]
 - Associates ACE metric (measure of cycle connectivity with the rest of the graph) for each cycle in graph
 - (d_{\max}, η) ACE constrained LDPC codes
 - Each cycle of length less than or equal to $2d_{\max}$ has ACE value larger or equal than η

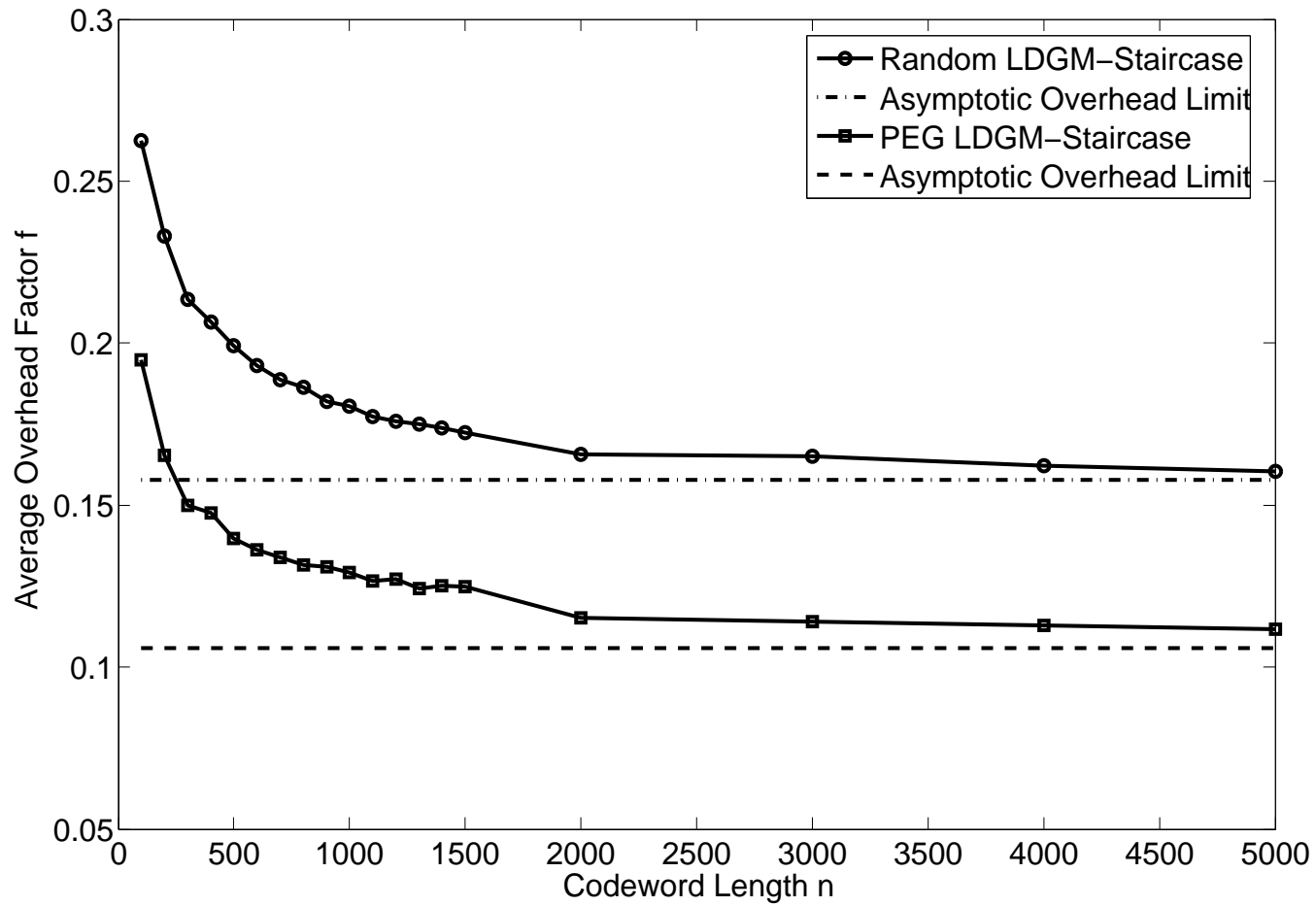


Improved LDGM–Staircase Design



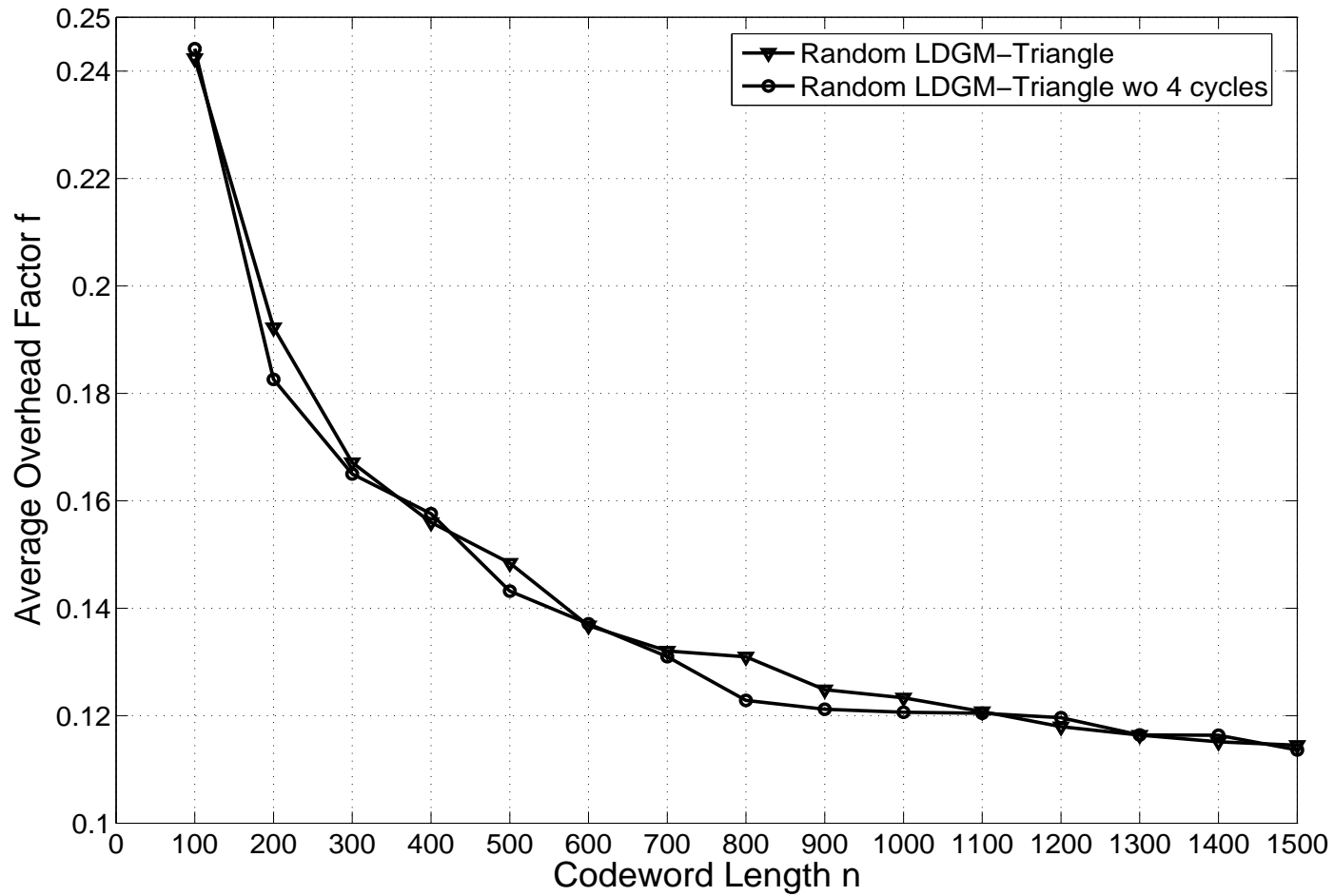


Asymptotic Overhead Factor





LDGM-Triangle Design





Conclusions and Future Work

- Overhead Factor
 - Performance measure of interest for Large Content Distribution Applications
 - Not much of a recent work
 - Patent-free constrained codes
 - Significant improvements over currently suggested codes are possible
- Future Work
 - Generally, extensive design analysis and simulation study is needed